

static games with incomplete information

part 1

roadmap

Bayes-Nash equilibrium (BNE)

examples

Harsanyi's proposal

definition

references

Chap. 20 of Dutta

example:

prisoners' dilemma - 2 is tough

1 type 1	2 Not to confess	Confess
	Not to confess	-1,-1
Confess	<u>0</u> , -6	<u>-3</u> , <u>-3</u>

example:

prisoners' dilemma - 2 is accommodating

1	2 type 2	Not to confess	Confess
		Not to confess	Confess
Not to confess		-1, <u>1</u>	-6,-2
Confess		<u>0</u> , <u>-4</u>	<u>-3</u> , <u>-5</u>

example:

coordination game: 2 is matched

1 type 1	2 Book launch	Movie
	Book launch	<u>2</u> , <u>1</u>
Movie	0, 0	<u>1</u> , <u>2</u>

example:

coordination game: 2 is mismatched

1 2 type 2	Book launch	Movie
	Book launch	<u>2</u> , 0
Movie	0, <u>2</u>	<u>1</u> , 0

static games with incomplete information

definition

a **game of incomplete information** is one in which players do not know some relevant characteristic of their opponents, which may include their payoffs, their available options, and even their beliefs

static games with incomplete information

examples

prisoners' dilemma

player 1 always plays Confess

player 2, type tough plays Confess

player 2, type accommodating plays Not
Confess

coordination game

player 1 plays mixed strategy λ

player 2, type matched plays mixed strategy μ_1

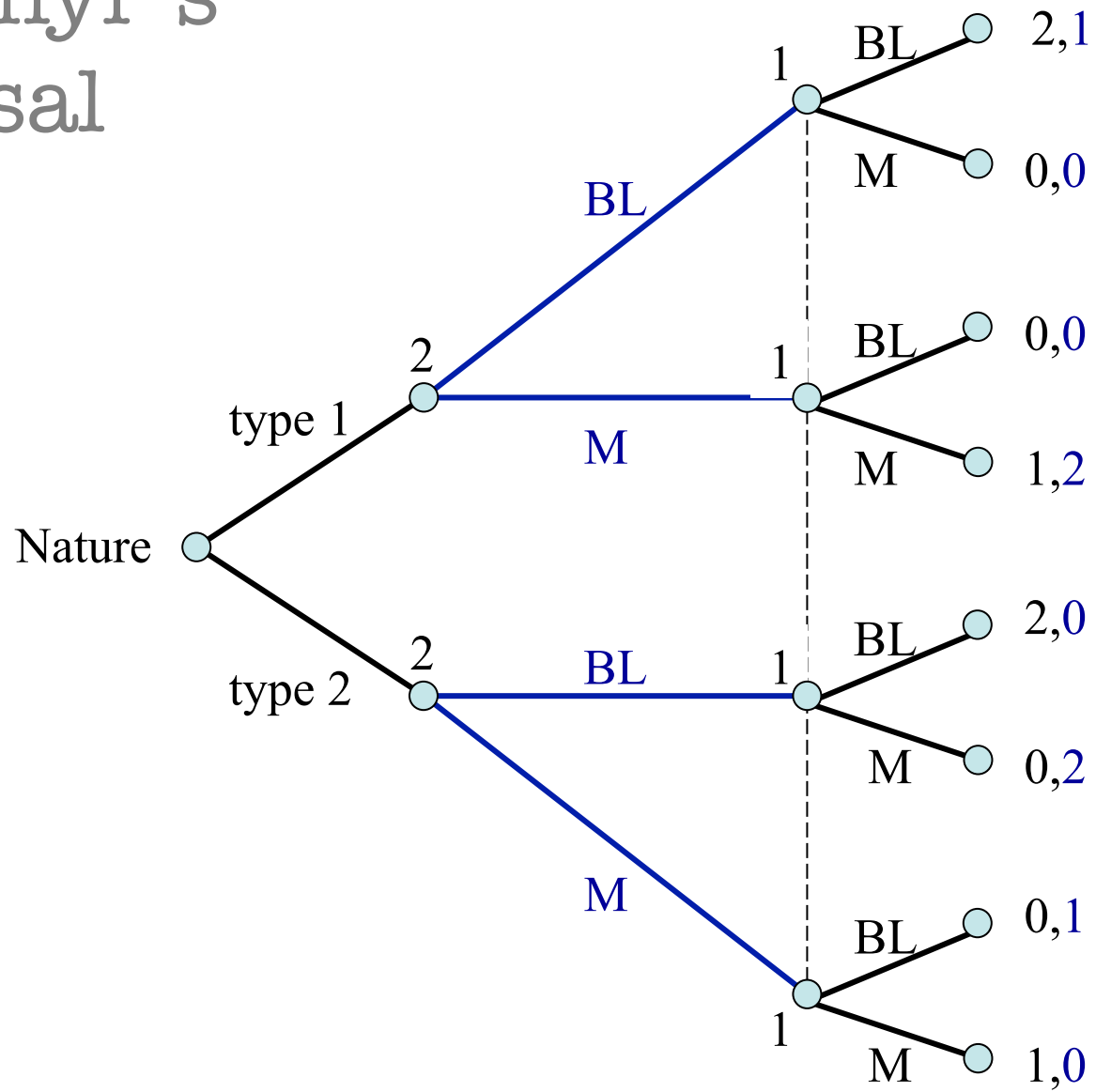
player 2, type mismatched plays mixed
strategy μ_2

static games with incomplete information
common priors

assumption of a common prior:

probabilities of types must become part of the
game and are known by all players

Harsanyi's proposal



static games with incomplete information

Bayes-Nash equilibrium

Harsanyi's proposal

- turn it into a game of complete but imperfect information
- use Nash equilibrium as the solution concept

a Bayes-Nash equilibrium of the game is a triple (λ, μ_1, μ_2) in which each player – and each player type – plays a best response, as follows:

- (1) μ_i maximizes the type i of player 2's payoff when λ is 1's strategy;
- (2) λ maximizes player 1's payoff when the type i of player 2 is playing μ_i and the probabilities of types 1 and 2 are respectively p and $1-p$

static games with incomplete information

Bayes-Nash equilibrium in the coordination game

two pure-strategy Bayes-Nash equilibria whenever $p \geq 2/3$:

1. player 1 plays BL while player 2 plays (BL,M)
2. player 1 plays M and player 2 plays (M,BL)

if $2/3 > p \geq 1/3$, there is only one pure-strategy Bayes-Nash equilibrium, namely the first one

if $p < 1/3$, there is no pure-strategy Bayes-Nash equilibrium

there is always a mixed-strategy Bayes-Nash equilibrium:

$$\lambda = 2/3, \mu_1 = 1/3, \mu_2 = 1/3$$

there may be additional Bayes-Nash equilibria,
characterized by $\lambda = 2/3$ and $p(3\mu_1 - 1) = (1-p)(1 - 3\mu_2)$

example

Cournot duopoly with incomplete information

A homogeneous product is produced by only two firms, 1 and 2. The quantities are denoted by q_1 and q_2

They choose the quantities simultaneously

The market price is given by $P(Q) = a - Q$, where $Q = q_1 + q_2$

Firm 1's cost function: $C_1(q_1) = c \cdot q_1$

All the above is common knowledge

example

Cournot duopoly with incomplete information

Firm 2's cost depends on some factor (e.g. technology) that only 2 knows. Its costs can be:

- High: $C_2(q_2) = c_H q_2$
- Low: $C_2(q_2) = c_L q_2$

Before production, firm 2 can observe the factor and know its marginal cost; however, firm 1 cannot observe firm 2's cost. She believes 2's cost function is

- High: $C_2(q_2) = c_H q_2$ with probability p
- Low: $C_2(q_2) = c_L q_2$ with probability $1-p$

example

Cournot duopoly with incomplete information

The BNE is $(q^*_2(c_H), q^*_2(c_L), q^*_1)$ such that:

$$q^*_2(c_H) = \frac{1}{3}(a - 2c_H + c) + \frac{1-p}{6}(c_H - c_L)$$

$$q^*_2(c_L) = \frac{1}{3}(a - 2c_L + c) + \frac{p}{6}(c_H - c_L)$$

$$q^*_1 = \frac{a - 2c + pc_H + (1-p)c_L}{3}$$

example:
coordination game

1 2	Book launch	Movie
	Book launch	$2+t_1, 1$
Movie	$0, 0$	$1, 2+t_2$

example:

coordination game - BNE

t_i follows $U [0,x]$

Bayes-Nash equilibrium:

player 1 plays BL if t_1 above c and M otherwise

player 2 plays M if t_2 above p and BL otherwise

player 1 plays BL with probability $(x-c)/x$

player 2 plays M with probability $(x-p)/x$

but $\lim_{x \rightarrow 0} (x-c)/x = \lim_{x \rightarrow 0} (x-p)/x = 2/3$